

## 一、 第一章

数据结构定义、数据定义:

### 1. Data

is the carrier of information.

Data is a set of **numbers** , **characters**, and **other symbols**

that can be used to describe the objective things.

These symbols can be input into computers , identified and processed by the computer program.

### 2 . Data structure

A data structure is a data object together with the relationships among the data members that compose the object

$\text{Data\_Structure} = \{D, R\}$

D is a data object,

R is a limited set of relationships of all the data members in D.

递归的概念与实现

例2. 求数组中的最大值

```
public static int findMax(int[] a, int n){
    //n表示n个元素，它们在数组a中
    if(n==1){
        return a [0];
    }
    else{
        int temp=findMax(a,n-1);
        return temp>a [n-1]?temp:a [n-1];
    }
}
```

```
int max(int a[],int n)
{ if(n == 1) return a[0];
  int m = max(a,n-1);
  if( m > a[n-1] )
    return m;
  else
    return a[n-1];
}
```

## 例5. 交换左右子树

```
void Swapchild ( BinTreeNode * p )
{ if ( p == NULL ) return ;
  BinTreeNode * temp = p -> left ;
  p ->left = p -> right ;
  p -> right = temp;
  Swapchild ( p ->left );
  Swapchild ( p ->right );
}
```

面向对象部分定义不会考

二、 第二章：算法分析

时间复杂度和空间复杂度

几个表示法，四种表示法

给一个算法计算对应的复杂度

计算明确的操作次数，会告知所有的相应操作的操作个数

最佳、最差和平均情况下的复杂度差异；

大 O、Ω和 θ 符号

1) 分析某个语句的执行次数（频度）

2) 分析某个程序段执行的时间复杂度（用大 O 表示，要求写出推导过程）

例 2. x = 0; y = 0;

```
for (int i = 1; i <= n; i++)
    for (int j = 1; j <= i; j++)
        for (int k = 1; k <= j; k++)
            x = x+y;
```

次数为:  $n*(n+1)*(n+2)/6$

例 3. int x = 91; int y = 100;

```
while(y>0)
{ if(x>100) { x -= 10; y--;}
  else x++;
}
```

1100 次

## 2.1 Space Complexity

2) example:

- Sequential Search

```
public static int SequentialSearch( int [ ] a , int x )
{
    int i;
    for(i=0; i<a.length && a[i]!=x; i++) ;
    if(i== a.length) return -1;
    return i;
}
```

## 2.1 Space Complexity

Total data space:

12 bytes : x,i,a[i],0,-1,a.length

栈空间

each of them cost 2 bytes

$S(n)=0$  常数复杂度

16位电脑  
在64位电脑上应为8 bytes

- Recursive code to add  $a[0:n-1]$

```
public static float Rsum(float [ ] a, int n)
{
    if ( n>0 )
        return Rsum(a, n-1) + a[n-1];
    return 0;
}
```

Recursion stack space:

formal parameters : a (2 byte), n(2 byte)

return address(2 byte)

Depth of recursion: n+1

$$S_{Rsum}(n) = 6(n+1)$$

一、第三章

线性表的定义

线性表的代码以及各种代码的变体

## ADT specification of a linear list

AbstractDataType LinearList

{ instances

ordered finite collections of zero or more elements

operations

Create(); Destroy();

IsEmpty(); Length();

Find(k,x); Search(x);

Delete(k,x); Insert(k,x);

Output(out);

}

## 2. Class definition

ListNode —— 代表结点的类

LinkedList —— 代表表本身的类

LinkedListItr —— 代表位置的类

} 都是包DataStructures的一部分

## 1) ListNode class



```
package DataStructures;
class ListNode
{   ListNode( object theElement)
    {   this( theElement, null);
        }
    ListNode( object theElement, ListNode n)
    {   element = theElement;
        next = n;
        }
    object element;
    ListNode next;
}
```

```
package DataStructures
public class LinkedListItr
{   LinkedListItr( ListNode theNode)
    {   current = theNode;
        }
    public boolean isPastEnd( )
    {   return current == null;
        }
    public object retrieve( )
    {   return isPastEnd( ) ? Null : current.element;
        }
    public void advance( )
    {   if( ! isPastEnd( ) )
        current = current.next;
        }
    ListNode current;
}
```

```

package DataStructures;
public class LinkedList
{
    public LinkedList()
        { header = new ListNode( null ) ; }
    public boolean isEmpty()
        { return header.next == null ; }
    public void makeEmpty()
        { header.next = null; }
    public LinkedListItr zeroth()
        { return new LinkedListItr( header ) ; }
    public LinkedListItr first()
        { return new LinkedListItr( header.next ) ; }
    public LinkedListItr find( object x )
    public void remove( object x )
    public LinkedListItr findPrevious( object x )
    public void insert( object x, LinkedListItr p )

    private ListNode header;
}

```

```

public LinkedListItr find (object x)
{
    ListNode itr = header.next;
    while ( itr != null && !itr.element.equals( x ) )
        itr = itr.next;
    return new LinkedListItr( itr );
}

```

O(N)

- Insert(x, p)

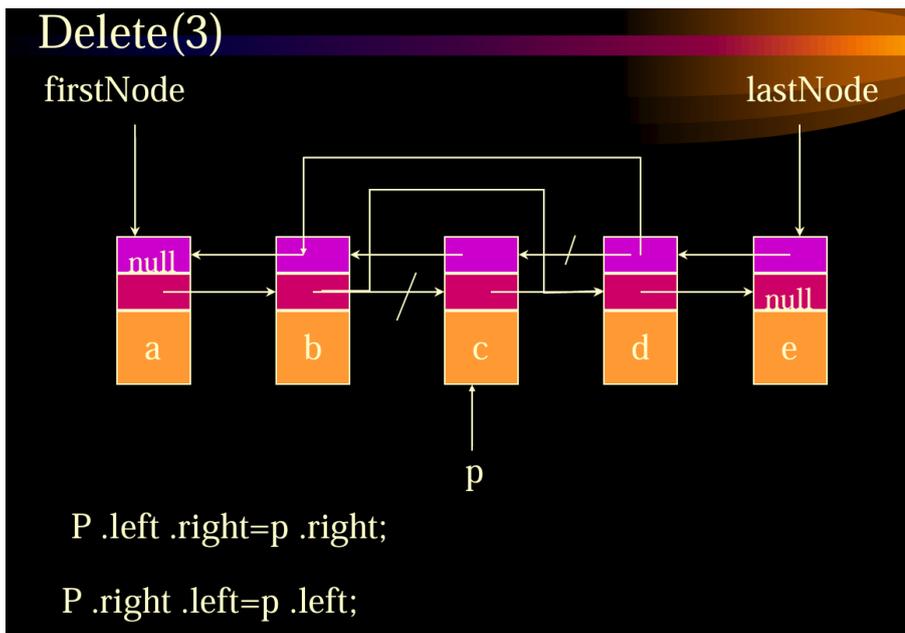
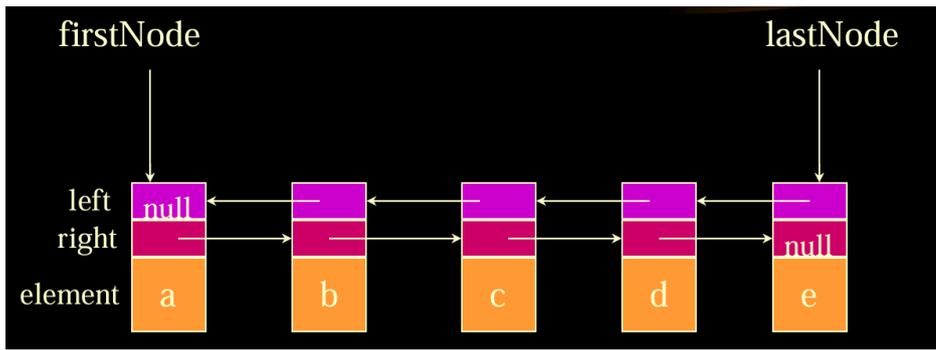
```

public void insert( object x, LinkedListItr p)
{
    if( p!=null && p.current != null )
        p.current.next = new ListNode( x, p.current.next );
}

```

O(1)

双向链表



静态链表和含有游标的链表

### 3.2.6. Cursor implementation of Linked Lists

use array to implement linked list:

cursorSpace		header	0	1	2	3	4	5	6	7	8
element	next		0	1	2	3	4	5	6	7	8
0	1			30	2						
1	2	p		50	3						
2	3			67	8						
3	4			15	7						
				81	0						
				10	4						
				20	1						
				78	5						

$P = p.next$        $p = cursorSpace[p].next$   
 $p.data$            $cursorSpace[p].data$

```

class CursorNode
{ CursorNode( object theElement )
    { this( theElement, 0 ); }
  CursorNode( object theElement, int n )
    { element = theElement;
      next = n;
    }

    object element;
    int next;
}

```

```

public class CursorListItr
{ CursorListItr( int theNode ) { current = theNode; }
  public boolean isPastEnd() { return current == 0; }
  public object retrieve()
  { return isPastEnd() ? null:
    CursorList.cursorSpace[ current ].element;
  }
  public void advance()
  { if( !isPastEnd() )
    current = CursorList.cursorSpace[ current ].next;
  }

  int current;
}

```

2) Class skeleton for CursorList

```

public class CursorList
{ private static int alloc()
  private static void free( int p)
  public CursorList()
    { header = alloc(); cursorSpace[ header ].next = 0; }
  public boolean isEmpty()
    { return cursorSpace[ header ].next == 0; }
  public void makeEmpty()
  public CursorListItr zeroth()
    { return new CursorListItr( header ); }
  public CursorListItr first()
    { return new CursorListItr( cursorSpace[ header ].next ); }
}

```

```

public CursorListItr find( object x )
public void insert( object x, CursorListItr p)
public void remove( object x )
public CursorListItr findPrevious( object x )

private int header;
static CursorNode [ ] cursorSpace;

private static final int SPACE-SIZE = 100;

static
{
  cursorSpace = new CursorNode[ SPACE-SIZE ];
  for( int i = 0; i<SPACE-SIZE; i++)
    cursorSpace[ i ] = new CursorNode( null, i + 1 );
  cursorSpace[ SPACE-SIZE-1].next = 0;
}
}

```

## Some Routines:

- Alloc and free

```

private static int alloc( )
{
  int p = cursorSpace[ 0 ].next;
  cursorSpace[0].next = cursorSpace[p].next;
  if( p == 0 )
    throw new OutOfMemoryError();
  return p;
}

private static void free( int p )
{
  cursorSpace[p].element = null;
  cursorSpace[p].next = cursorSpace[0].next;
  cursorSpace[0].next = p;
}

```

多项式相加:

\*问题:  $A(X) = 2X^{100} + 3X^{14} + 2X^8 + 1$

$$B(X) = -2X^{100} + 8X^{14} - 3X^{10} + 10X^6 - X$$

$$A(X) + B(X) = 11X^{14} - 3X^{10} + 2X^8 + 10X^6 - X + 1$$

方法: 设4个引用变量:

pa, pb, pc, p(c++需要)

1) 初始化: pc, pa, pb;

2) 当pa和pb都有项时

pc永远指向相加时结果链表的最后一个结点。

a) 指数相等 ( pa.exp == pb.exp )

对应系数相加: pa.coef = pa.coef + pb.coef;

p = pb(c++需要); pb前进;

if (系数相加结果为0) { p = pa; pa前进; }

else { pc.link = pa; pc = pa; pa前进 }

b) 指数不等 pa.exp < pb.exp //pb要插入结果链表

{ pc.link = pb; pc = pb; pb前进 }

c) 指数不等 pa.exp > pb.exp //pa要插入结果链表

{ pc.link = pa; pc = pa; pa前进 }

3) 当两链表中有一链表为空, 则将另一链表链入结果链表就可以

if (pb空了) { pc.link = pa; }

else pc.link = pb;

## Linked List Implementation of Stacks

```
public class StackLi
{
    public StackLi() { topOfStack = null; }
    public boolean isFull() { return false; }
    public boolean isEmpty() { return topOfStack == null; }
    public void makeEmpty() { topOfStack = null; }

    public void push( object x )
    public object top()
    public void pop() throws Underflow
    public object topAndPop()

    private ListNode topOfStack;
}
```

## Array Implementation of Stacks

```
public class stackAr
{
    public StackAr()
    public StackAr( int capacity )

    public boolean isEmpty() { return topOfStack == -1; }
    public boolean isFull() { return topOfStack == theArray.length - 1; }
    public void makeEmpty() { topOfStack = -1; }

    public void push( object x ) throws overflow
    public object top()
    public void pop() throws Underflow
    public object topAndPop()

    private object [ ] theArray;
    private int topOfStack;

    static final int DEFAULT_CAPACITY = 10;
}
```

栈、队列的例子（重点看，括号匹配、表达式计算）

```
#include <iostream.h>
#include <string.h>
#include <stdio.h>
#include "stack.h"
const int Maxlength = 100; // max expression length
void PrintMatchedPairs(char *expr)
{ Stack<int> s(Maxlength);
  int j, length = strlen(expr);
  for ( int i = 1; i <= length; i++)
  { if ( expr[i-1] == '(') s.Add(i);
    else if (expr[i-1] == ')')
      try {s.Delete(j); cout <<j<<" " <<i<< endl;}
        catch (OutOfBounds)
          {cout << "No match for right parenthesis"
            << " at " << i << endl;}
    }
  while ( !s.IsEmpty ()
  { s.Delete(j);
    cout<< "No match for left parenthesis at "
      << j << endl;
  }}
}
```

## 3.4 . The Queue ADT

A queue is a linear list in which additions and deletions take place at different ends.

It is also called a first-in-first-out list.

The end at which new elements are added is called the **rear**.

The end from which old elements are deleted is called the **front**.

AbstractDataType Queue

```
{
  instances
    ordered list of elements;one end is called the front; the other is the rear;
  operations
    Create(): Create an empty queue;
    IsEmpty(): Return true if queue is empty,return false otherwise;
    IsFull(): return true if queue is full, return false otherwise;
    First(): return first element of the queue;
    Last(): return last element of the queue;
    Add(x): add element x to the queue;
    Delete(x): delete front element from the queue and put it in x;
}
```

杨晖三角 (不用重点看)